

**INVENTION:** Multiple Data Rate Complex Walsh Codes for CDMA

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5 **TECHNICAL FIELD**

The present invention relates to CDMA (Code Division Multiple Access) cellular telephone and wireless data communications with data rates up to multiple T1 (1.544 Mbps) and higher (>100 Mbps), and to optical CDMA with data rates in the  
10 Gbps and higher ranges. Applications are mobile, point-to-point and satellite communication networks. More specifically the present invention relates to novel multiple data rate algorithms for complex and hybrid complex Walsh orthogonal CDMA codes. These algorithms generate multiple code length complex Walsh and  
15 hybrid complex Walsh orthogonal codes for use as the channelization codes for multiple data rate users. These new algorithms and codes offer substantial improvements over the current real Walsh orthogonal variable spreading factor (OVSF) CDMA codes for the next generation wideband CDMA (W-CDMA).

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30 **ABSTRACT**

The present invention describes new multiple data rate algorithms for complex Walsh and hybrid complex Walsh orthogonal CDMA channelization encoding and decoding of multiple data rate  
35 users. Complex Walsh and hybrid complex Walsh orthogonal CDMA

codes have been disclosed in a previous patent application [6]  
for constant data rate communications. These algorithms  
generate a means to accomodate multiple data rate users over the  
same CDMA frequency band using complex Walsh and hybrid complex  
40 Walsh orthogonal codes.

The means of this invention is to provide complex Walsh and  
hybrid complex Walsh with the means to separate the different  
data rate users in the sequency domain of the complex Walsh  
analogous to the current use of different frequency bands for the  
45 different data rate users. Sequency for complex Walsh and hybrid  
complex Walsh codes is the average rate of phase angle rotations  
of the code vectors, and is analogous to frequency in the Fourier  
domain. For example, a frequency equal to 1 rotation per N chip  
code length corresponds to a sequency equal to 1 complete phase  
50 rotation over the N chip code vector length. However, for real  
Walsh codes the definition of sequency is the average number of  
sign changes of the code vectors, and has no direct analogy to  
frequency. The widest sequency bands consisting of sets of  $N/2$   
contiguous sequencies is reserved for the highest data rate users  
55 who transmit at a 2 chip code rate, the next widest sequency  
bands consisting of  $N/4$  contiguous sequencies is reserved for  
the 4 chip code rate users,..., and the narrowest sequeny bands  
consisting of a single sequency are reserved for the N chip code  
rate users.

60 Current art uses algorithms to generate multiple code  
length real Walsh CDMA orthogonal codes for the next generation  
wideband CDMA (W-CDMA). These codes are orthogonal variable  
spreading factor (OVSF) CDMA codes. Spreading factor is the  
length of each code, and a variable spreading factors refers to a  
65 variable code length or equivalently to a choice of multiple  
code lengths.

The present invention provides a means to significantly  
improve CDMA performance for multiple data rate users by allowing  
the use of the new complex Walsh and hybrid complex Walsh CDMA

orthogonal codes in place of the real Walsh OVSF CDMA orthogonal codes and with implementation means for fast and computationally efficient encoding and decoding. Fundamental improvements are known to include an increase in the carrier-to-noise ratio (CNR) for data symbol recovery in the receiver, lower correlation side-lobes under timing offsets both with and without PN spreading, lower levels of harmonic interference caused by non-linear amplification of multi-carrier CDMA signals, and reduced phase tracking jitter for code tracking to support both acquisition and synchronization. Hybrid complex Walsh orthogonal CDMA codes increase the choices for the code lengths by allowing the combined use of complex Walsh, Hadamard, and discrete Fourier transform complex orthogonal codes using a Kronecker construction, direct sum construction, as well as the possibility for more general functional combining.

#### BACKGROUND ART

Current art is represented by the work on orthogonal spreading factor (OVSF) real Walsh codes for wideband CDMA (W-CDMA) for the third generation CDMA (G3) proposed standard candidates and for broadband wireless communications, and the previous work on the real Walsh fast transform algorithms. These are documented in references 1,2,3,4,5. Reference 1 is an issue of the IEEE communications journal devoted to wideband CDMA including OVSF. References 2 and 3 are issues of the IEEE communications magazine that are devoted to "Multiple Access for Broadband Networks" and "Wideband CDMA". Reference 4 is an issue of the IEEE personal communications devoted to "Third Generation Mobile Systems in Europe". Reference 5 is the widely used reference on real Walsh technology which includes algorithms for the fast Walsh transform. The new complex Walsh and hybrid complex Walsh orthogonal CDMA codes being addressed in this invention for application to multiple data rate users, have been disclosed in a previous patent application [6] for constant data rate communications.

Current art using real Walsh orthogonal CDMA channelization codes to generate OVSF codes for multiple data rate users is represented by the scenario described in the following with the aid of equations (1) and (2) and FIG. 1,2,3,4. This scenario considers CDMA communications spread over a common frequency band for each of the communication channels. These CDMA communications channels for each of the multiple rate users are defined by assigning a unique real Walsh orthogonal spreading code to each user. This real Walsh code has a maximum length of N chips with  $N=2^M$  where M is an integer, with shorter lengths of 2,4,...,N/2 for the higher data rate users. These multiple length real Walsh codes have limited orthogonality properties and occupy the same frequency band. These Walsh encoded user signals are summed and then re-spread over the same frequency band by PN codes, to generate the CDMA communications signal which is modulated and transmitted. The communications link consists of a transmitter, propagation path, and receiver, as well as interfaces and control.

It is assumed that the communication link is in the communications mode with the users communicating at symbol rates equal to the code repetition rates of their respective communications channels and that the synchronization is sufficiently accurate and robust to support this communications mode. In addition, the power differences between users due to differences in data rates and in communication link budget parameters is assumed to be incorporated in the data symbol amplitudes prior to the CDMA encoding in the CDMA transmitter, and the power is uniformly spread over the wideband by proper selection of the CDMA pulse waveform. It is self evident to anyone skilled in the CDMA communications art that these communications mode assumptions are both reasonable and representative of the current CDMA art and do not limit the applicability of this invention.

**Transmitter equations** (1) describe a representative real Walsh CDMA encoding for multiple data rate users for the

transmitter in FIG. 1. These equations represent a considerably more sophisticated and improved implementation of current OVSF CDMA communications which has been developed to help support the new invention for complex Walsh and hybrid complex Walsh CDMA orthogonal codes.

Lowest data rate users are assumed to communicate at the lowest symbol rate equal to the code repetition rate of the  $N$  chip real Walsh code, which means they are assigned  $N$  chip code vectors from the  $N \times N$  real Walsh code matrix  $W_N$  in 1 for their channelization codes. Higher data rate users will use shorter real Walsh codes. The reference real Walsh code matrix  $W_N$  has  $N$  Walsh row code vectors  $W_N(c)$  each of length  $N$  chips and indexed by  $c=0,1,\dots,N-1$ , with  $W_N(c)=[W_N(c,1),\dots,W_N(c,N)]$  wherein  $W_N(c,n)$  is chip  $n$  of code  $u$ . Walsh code chip  $n$  of code vector  $u$  has the possible values  $W_N(c,n)=+/-1$ .

Multiple data rate menu in 2 lists the possible user data symbol rates  $R_s$  and the corresponding code lengths and symbols transmitted over each  $N$  chip reference code length. User symbol rate  $R_s=1/NT$  is the code repetition rate  $1/NT$  of the  $N$  chip code over the code time interval  $NT$ . User data rate  $R_b$  in bits/second is equal to  $R_b=R_s b_s$  where  $b_s$  is the number of data bits encoded in each data symbol. Assuming a constant  $b_s$  for all of the multiple data rate users, the user data rate becomes directly proportional to the user symbol rate  $R_b \sim R_s$  which means the user symbol rate menu in 1 is equivalent to the user data rate menu.

User data symbols and channelization codes are listed in 3 for the multiple rate users. Users are grouped into the data rate categories corresponding to their respective code chip lengths  $2,4,8,\dots,N/2,N$  chips. User groups are indexed by  $m=1,2,\dots,M$  where group  $m$  consists of all users with  $N(m)=2^m$  chip length codes drawn from the  $N(m) \times N(m)$  real Walsh code matrix  $W_{N(m)}$ . Users within group  $m$  are identified by the index  $u_m$  which is set equal to the Walsh channelization code vector index in  $W_{N(m)}$ . Code chip  $n_m$  of the user code  $u_m$  is equal to  $W_{N(m)}(u_m, n_m)$

175 where  $n_m=0,1,2,\dots,N(m)-1$  is the chip index. User data symbols  
 $Z(u_{m,k_m})$  are indexed by  $u_{m,k_m}$  where the index  $k_m=0,1,2,\dots,N/N(m)-1$   
 identifies the data symbols of  $u_m$  which are transmitted over the  
 N chip code block. The total number of user data symbols  
 transmitted per N chip block is N which means the number of  
 180 channel assignments  $\{u_m, m=1,2,\dots,M\}$  will be less than N for  
 multiple data rate CDMA communications when there is at least one  
 user using a higher data rate.

**Current multiple data rate real Walsh CDMA encoding (1)  
 for transmitter**

185 **1 N chip Walsh code block**

$W_N$  = Walsh NxN orthogonal code matrix consisting of  
 N rows of N chip code vectors  
 =  $[W_N(c)]$  matrix of row vectors  $W_N(c)$   
 =  $[W_N(c,n)]$  matrix of elements  $W_N(c,n)$   
 190  $W_N(c)$  = Walsh code vector c for  $c=0,1,\dots,N-1$   
 =  $[W_N(c,0), W_N(c,1), \dots, W_N(c,N-1)]$   
 = 1xN row vector of chips  $W_N(c,0), \dots, W_N(c,N-1)$   
 $W_N(c,n)$  = Walsh code c chip n  
 = +/-1 possible values

195 **2 Multiple data rate menu**

N chip real Walsh symbol rate

$R_s$  = User symbol rate, symbols/second  
 =  $1/NT$  where T = Chip repetition interval

Symbol rate menu for multiple data rates

	Symbol rate, Symbols/second	Code length, chips	Symbols per N chips
$R_s$ =	$1/2T$	2	$N/2$
=	$1/4T$	4	$N/4$
=	$1/8T$	8	$N/8$
	$\vdots$	$\vdots$	$\vdots$
=	$1/2NT$	$N/2$	2
	$= 1/NT$	N	1

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### 3 User data symbols and channelization codes

Users are categorized into  $M$  groups according to the number of code chips.

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$m$  = Index of the user groups  
 $= 1, 2, \dots, M$

$u_m$  = One of up to  $N(m)=2^m$  possible users in group  $m$

$N(m)$  = Number of code chips for the codes in the user group  $m$   
 $= 2^{m+1}$

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User data symbols

$Z(u_{m,k_m})$  = User  $u_m$  data symbol  $k_m$

$k_m$  = Index for the user data symbols over the  $N$  chip code block, for a user from group  $m$   
 $= 0, 1, 2, \dots, N/N(m)-1$

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User channelization codes within each group are selected from a subset of the orthogonal codes in the Walsh code matrix.

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$W_{N(m)}(u_m)$  = Walsh  $1 \times 2^m$  dimensional code vector  $u_m$  in the  $N(m) \times N(m)$  Walsh code matrix, for user  $u_m$  in the group  $m$

$W_{N(m)}(u_m, n_m)$  = User  $u_m$  code chip  $n_m=0, 1, 2, \dots, N(m)-1$

### 4 Real Walsh encoding and channel combining

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$\tilde{Z}(n)$  = Real Walsh CDMA encoded chip  $n$

$$= \sum_{m=1}^M \sum_{u_m} Z(u_{m,k}) W_{N(m)}(u_m, n=n_m + k_m N(m))$$

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## 5 PN scrambling

$P_R(n), P_I(n)$  = PN code chip  $n$  for real, imaginary axes  
axes

$Z(n)$  = PN scrambled real Walsh encoded data chips  
after summing over the users

$$= \sum_u \tilde{Z}(n) [P_R(n) + jP_I(n)]$$

$$= \sum_u \tilde{Z}(n) [\text{sgn}\{P_R(n)\} + j\text{sgn}\{P_I(n)\}]$$

= Real Walsh CDMA encoded complex chips  
after PN scrambling

Walsh encoding and channel combining in 4 encodes each of the users  $\{u_m\}$  and their data symbols  $\{Z(u_{m,k_m})\}$  with a Walsh code  $W_{N(m)}(u_m)$  drawn from the group  $m$  of the  $N(m)$  chip channelization codes where  $u_m$  is the user code. A time delay of  $k_m N(m)$  chips before start of the real Walsh encoding of the data symbol  $k_m$  in each of the user channels, is required for implementation of the multiple data rate user real Walsh encoding and for the summation of the encoded data chips over the users. Output of this multiple data rate real Walsh encoding and summation over the multiple data rate users is the set of real

Walsh CDMA encoded chips  $\{\tilde{Z}(n)\}$  over the  $N$  chip block.

PN scrambling of the real Walsh CDMA encoded chips in 5 is accomplished by encoding the  $\{\tilde{Z}(n)\}$  with a complex PN which is constructed as the complex code sequence  $[P_R(n) + jP_I(n)]$  wherein  $P_R(n)$  and  $P_I(n)$  are independent PN sequences used for the real and imaginary axes of the complex PN. These PN codes are 2-phase with each chip equal to  $\pm 1$  which means PN encoding consists of sign changes with each sign change corresponding to the sign of the PN chip. Encoding with PN means each chip of the summed Walsh encoded data symbols has a sign change when the corresponding PN chip is  $-1$ , and remains unchanged for  $+1$  values. This operation is described by a multiplication of each chip of the summed Walsh encoded data symbols with the sign of



the PN chip. Purpose of the PN encoding for complex data symbols is to provide scrambling of the summed Walsh encoded data symbols as well as isolation between groups of users. Output of this real Walsh CDMA encoding followed by the complex PN scrambling are the CDMA encoded chips over the N chip block  $\{Z(n)\}$ .

**Receiver equations (2)** describe a representative multiple data rate real Walsh CDMA decoding for the receiver in FIG. 3.

The receiver front end **5** provides estimates  $\{\hat{Z}(n)\}$  of the transmitted real Walsh CDMA encoded chips  $\{Z(n)\}$ . Orthogonality property **6** is expressed as a matrix product of the real Walsh code chips or equivalently as a matrix product of the Walsh code chip numerical signs, for any of the 2,4,8,...,N/2,N chip real Walsh channelization codes and their repetitions over the N chip code block. These codes are orthogonal with respect to the user codes within a group. They are also orthogonal between code groups for the allowable subsets of code assignments to the users, for all code repetitions over the N chip code block. This means that the allowable codes  $\{u_m\}$  in group m are orthogonal to the allowable codes  $\{u_{m+p}\}$  in group m+p for all code repetitions of the codes  $\{u_m\}$  over the N chip code block, for  $p \geq 0$ .

The 2-phase PN codes **7** have the useful decoding property that the square of each code chip is unity which is equivalent to observing that the square of each code chip numerical sign is unity. Decoding algorithms **8** perform the inverse of the signal processing for the encoding in equations (1) to recover estimates  $\{\hat{Z}(u_{m,k_m})\}$  of the transmitter user symbols  $\{Z(u_{m,k_m})\}$ .

**Current multiple data rate real Walsh CDMA decoding (2)  
for receiver**

**5** Receiver front end provides estimates  $\{\hat{Z}(n) = \hat{R}(n) + j\hat{I}(n)\}$  of the encoded transmitter chip symbols  $\{Z(n)\}$

**6** Orthogonality properties of the set of real Walsh  $\{2 \times 2, 4 \times 4, 8 \times 8, \dots N \times N\}$  matrices

The  $N(m) \times N(m)$  Walsh code matrices for all  $m$  are orthogonal

$$N(m)^{-1} \sum_{n_m} W_{N(m)}(\hat{c}_m, n_m) W_{N(m)}(n_m, c_m) = \delta(\hat{c}_m, c_m)$$

where  $c_m, n_m = 0, 1, \dots, N(m)$

$$\begin{aligned} \delta(\hat{c}_m, c_m) &= \text{Delta function of } \hat{c}_m \text{ and } c_m \\ &= 1, 0 \quad \text{for } \hat{c}_m = c_m, \text{ otherwise} \end{aligned}$$

The  $N(m) \times N(m)$  and  $N(m+p) \times N(m+p)$  Walsh code matrices for all  $m$  and  $p \geq 0$  are orthogonal for a subset of codes  $\{u_m\}$  and  $\{u_{m+p}\}$

$$\begin{aligned} N(m)^{-1} \sum_{n_m} W_{N(m)}(u_m, n_m) \bullet W_{N(m+p)}(u_{m+p}, n_{m+p} = n_m + k_m N(m)) \\ = 0 \quad \text{for } k_m = 0, 1, 2, \dots, N/N(m) - 1 \end{aligned}$$

#### 7 PN decoding property

$$\begin{aligned} P(n) P(n) &= \text{sgn}\{P(n)\} \text{sgn}\{P(n)\} \\ &= 1 \end{aligned}$$

#### 8 Decoding algorithm

$$\begin{aligned} \hat{Z}(u_{m,k_m}) &= \\ N^{-1} \sum_{n_m} \hat{Z}(n) [\text{sgn}\{P_R(n)\} - j \text{sgn}\{P_I(n)\}] \bullet \\ \text{sgn}\{W_{N(m)}(n = n_m + k_m N(m), u_m)\} \\ &= \text{Receiver estimate of the transmitted complex} \\ &\quad \text{data symbol } Z(u_{m,k_m}) \end{aligned}$$

**FIG. 1 CDMA transmitter block diagram** is representative of a current CDMA transmitter which includes an implementation of the current multiple data rate real Walsh CDMA channelization encoding in equations (1). This block diagram becomes a representative implementation of the CDMA transmitter which implements the new multiple data rate complex Walsh and hybrid complex Walsh CDMA encoding, when the current multiple data rate real Walsh CDMA encoding **13** is replaced by the new

335 multiple data rate complex Walsh and hybrid complex Walsh CDMA  
encoding of this invention.

Signal processing starts with the stream of user input data  
words **9**. Frame processor **10** accepts these data words and  
performs the encoding and frame formatting, and passes the  
340 outputs to the symbol encoder **11** which encodes the frame  
symbols into amplitude and phase coded symbols  $\{Z(u_{m,k})\}$  **12**.  
These symbols **12** are the inputs to the current multiple data rate  
real Walsh CDMA encoding in equations (1). Inputs  $\{Z(u_{m,k})\}$  **12**  
are real Walsh encoded, summed over the users, and scrambled by  
345 complex PN in the current multiple data rate real Walsh CDMA  
encoder **13** to generate the complex output chips  $\{Z(n)\}$  **14**. This  
encoding **13** is a representative implementation of equations (1).  
These output chips  $Z(n)$  are waveform modulated **15** to generate  
the analog complex signal  $z(t)$  which is single sideband  
350 upconverted, amplified, and transmitted (Tx) by the analog front  
end of the transmitter **15** as the real waveform  $v(t)$  **16** at  
the carrier frequency  $f_0$  whose amplitude is the real part of the  
complex envelope of the baseband waveform  $z(t)$  multiplied by the  
carrier frequency and the phase angle  $\phi$  accounts for the phase  
355 change from the baseband signal to the transmitted signal.

It should be obvious to anyone skilled in the  
communications art that this example implementation in FIG. 1  
clearly defines the fundamental CDMA signal processing relevant  
to this invention disclosure and it is obvious that this example  
360 is representative of the other possible signal processing  
approaches.

**FIG. 2 multiple data rate real Walsh CDMA encoding** is a  
representative implementation of the multiple data rate real  
Walsh CDMA encoding **13** in FIG. 1 and in equations (1). Inputs  
365 are the complex user data symbols  $\{Z(u)\}$  **17**. Encoding of each  
user by the corresponding Walsh code is described in **18** by the  
implementation of transferring the sign of each Walsh code chip  
to the user data symbol followed by a 1-to-N expander  $1 \uparrow N$  of each

data symbol into an N chip sequence using the sign transfer of  
 370 the Walsh chips. The sign-expander operation **18** generates the  
 N-chip sequence  $Z(u_{m,k_m}) \text{sgn}\{W(u_m, (n=n_m+k_mN(m)))\}$  for  $n=0,1,\dots,N-1$   
 for each user  $\{u_m\}$ . This Walsh encoding serves to spread each  
 user data symbol into an orthogonally encoded chip sequence which  
 is spread over the CDMA communications frequency band. The Walsh  
 375 encoded chip sequences for each of the user data symbols are  
 summed over the users **19** followed by PN encoding with the  
 scrambling sequence  $[P_R(n)+jP_I(n)]$  **20**. PN encoding is  
 implemented by transferring the sign of each PN chip to the  
 summed chip of the Walsh encoded data symbols. Output is the  
 380 stream of complex multiple data rate real Walsh CDMA encoded  
 chips  $\{Z(n)\}$  **21**.

It should be obvious to anyone skilled in the  
 communications art that this example implementation in FIG. 2  
 clearly defines the fundamental CDMA signal processing relevant  
 385 to this invention disclosure and it is obvious that this example  
 is representative of the other possible signal processing  
 approaches.

**FIG. 3 CDMA receiver block diagram** is representative of a  
 current CDMA receiver which includes an implementation of the  
 390 current multiple data rate real Walsh CDMA decoding in equations  
 (2). This block diagram becomes a representative implementation  
 of the CDMA receiver which implements the new multiple data rate  
 complex Walsh and hybrid complex Walsh CDMA decoding when the  
 current multiple data rate real Walsh CDMA decoding **27** is  
 395 replaced by the new multiple data rate complex Walsh and hybrid  
 complex Walsh CDMA decoding of this invention.

FIG. 3 signal processing starts with the user transmitted  
 wavefronts incident at the receiver antenna **22** for the users  
 $\{u_m\}$ . These wavefronts are combined by addition in the antenna  
 400 to form the receive (Rx) signal  $\hat{v}(t)$  at the antenna output **22**  
 where  $\hat{v}(t)$  is an estimate of the transmitted signal  $v(t)$  **16** in  
 FIG. 1, that is received with errors in time  $\Delta t$ , frequency  $\Delta f$ ,

phase  $\Delta\theta$ , and with an estimate  $\hat{z}(t)$  of the transmitted complex baseband signal  $z(t)$  16 in FIG. 1. This received signal  $\hat{v}(t)$  is amplified and downconverted by the analog front end 23 and then synchronized and analog-to-digital (ADC) converted 24. Outputs from the ADC are filtered and chip detected 25 by the fullband chip detector, to recover estimates  $\{\hat{Z}(n)\}$  26 of the transmitted signal which is the stream of complex CDMA encoded chips  $\{Z(n)$  14 in FIG. 1. CDMA decoder 27 implements the algorithms in equations (2) by stripping off the PN code(s) and decoding the received CDMA real Walsh orthogonally encoded chips to recover estimates  $\{\hat{Z}(u_{m,k_m})\}$  28 of the transmitted user data symbols  $\{Z(u_{m,k_m})\}$  12 in FIG. 1. These estimates 28 are processed by the symbol decoder 29 and the frame processor 30 to recover estimates 31 of the transmitted user data words.

It should be obvious to anyone skilled in the communications art that this example implementation clearly defines the fundamental current CDMA signal processing relevant to this invention disclosure and it is obvious that this example is representative of the other possible signal processing approaches.

**FIG. 4 multiple data rate real Walsh CDMA decoding** is a representative implementation of the multiple data rate real Walsh CDMA decoding 27 in FIG. 3 and in equations (2). Inputs are the received estimates of the multiple data rate complex real Walsh CDMA encoded chips  $\{\hat{Z}(n)\}$  32. The PN scrambling code is stripped off from these chips 33 by changing the sign of each chip according to the numerical sign of the real and imaginary components of the complex conjugate of the PN code as per the decoding algorithms 8 in equations (2). Real Walsh channelization coding is removed in 34 by a pulse compression operation consisting of multiplying each received chip by the numerical sign of the corresponding Walsh chip for the user and

435 summing the products over the N Walsh chips to recover estimates  
 $\{\hat{Z}(u_{m,k_m})\}$  35 of the user complex data symbols  $\{Z(u_{m,k_m})\}$ .

It should be obvious to anyone skilled in the communications art that this example implementation clearly defines the fundamental current CDMA signal processing relevant  
440 to this invention disclosure and it is obvious that this example is representative of the other possible signal processing approaches.

For cellular applications the transmitter description describes the transmission signal processing applicable to this  
445 invention for both the hub and user terminals, and the receiver describes the corresponding receiving signal processing for the hub and user terminals for applicability to this invention.

For optical communications applications the the microwave processing at the front end of both the transmitter and the  
450 receiver is replaced by the optical processing which performs the complex modulation for the optical laser transmission in the transmitter and which performs the optical laser receiving function of the microwave processing to recover the complex baseband received signal.

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#### **SUMMARY OF INVENTION**

This invention is a new set of fast and computationally efficient algorithms for new multiple data rate orthogonal channelization encoding and decoding for CDMA using the new  
460 complex Walsh codes and the hybrid complex Walsh orthogonal codes in place of the current real Walsh orthogonal codes. Real Walsh codes are used for current CDMA applications and will be used for all of the future CDMA systems. The newly invented complex Walsh codes disclosed in [6] provide the choice of using the new  
465 complex Walsh codes or the real Walsh codes since the real Walsh codes are the real components of the complex Walsh codes. This means an application capable of using the complex Walsh codes can

simply turn-off the complex axis components of the complex Walsh codes for real Walsh CDMA coding and decoding.

470 Performance is improved for the multiple data rate CDMA communications when the new 4-phase complex Walsh orthogonal CDMA codes replace the current 2-phase real Walsh codes. These improvements include an increase in the carrier-to-noise ratio (CNR) for data symbol recovery in the receiver, lower  
475 correlation side-lobes under timing offsets both with and without PN spreading, lower levels of harmonic interference caused by non-linear amplification of multi-carrier CDMA signals, and reduced phase tracking jitter for code tracking to support both acquisition and synchronization. These potential performance  
480 improvements simply reflect the widely known principle that complex CDMA is better than real CDMA.

In addition to the performance improvement, there are greater code length choices for multiple data rate CDMA communications using the new hybrid complex Walsh orthogonal CDMA  
485 codes in place of the complex Walsh orthogonal CDMA codes which have been disclosed in [6]. Code length choices are increased by the combined use of complex Walsh and discrete Fourier transform complex orthogonal codes using a Kronecker construction, direct sum construction, as well as the  
490 possibility for more general functional combining.

#### BRIEF DESCRIPTION OF DRAWINGS

The above-mentioned and other features, objects, design algorithms, and performance advantages of the present invention  
495 will become more apparent from the detailed description set forth below when taken in conjunction with the drawings and performance data wherein like reference characters and numerals denote like elements, and in which:

FIG. 1 is a representative CDMA transmitter signal  
500 processing implementation block diagram with emphasis on the current multiple data rate real Walsh CDMA encoding which

contains the signal processing elements addressed by this invention disclosure.

505 FIG. 2 is a representative real Walsh CDMA encoding implementation diagram with emphasis on the current multiple data rate real Walsh CDMA encoding which contains the signal processing elements addressed by this invention disclosure.

510 FIG. 3 is a representative CDMA receiver signal processing implementation block diagram with emphasis on the current multiple data rate real Walsh CDMA decoding which contains the signal processing elements addressed by this invention disclosure.

515 FIG. 4 is a representative CDMA decoding implementation diagram with emphasis on the current multiple data rate real Walsh CDMA decoding which contains the signal processing elements addressed by this invention disclosure.

520 FIG. 5 is a representative CDMA encoding implementation diagram which describes the new complex Walsh and hybrid complex Walsh CDMA encoding of multiple data rate users and which contains the signal processing elements addressed by this invention disclosure.

525 FIG. 6 is a representative CDMA decoding implementation diagram which describes the new complex Walsh and hybrid complex Walsh CDMA decoding of multiple data rate users and which contains the signal processing elements addressed by this invention disclosure.

#### DISCLOSURE OF INVENTION

530 The new invention provides the algorithms and implementation architectures to support simultaneous multiple data rates or equivalently simultaneous multiple symbol rates using the new complex Walsh and hybrid complex Walsh orthogonal CDMA codes which have been disclosed in the invention application [6]. Simultaneous multiple data rates over the same CDMA  
535 frequency spectrum are well known in CDMA networking and been included in the next generation UMTS 3G evolving CDMA using



wideband CDMA (W-CDMA) and real Walsh orthogonal CDMA channelization codes.

The current art uses three categories of techniques designed to accommodate multiple data rate users and these are A) multiple chip length codes for the multiple data rate users, B) same chip length codes with the number of codes adjusted as required for the multiple data rate users, and C) different frequency spectrums assigned to the multiple data rate users which is frequency division multiplexing (FDM). The first technique is the preferred choice for W-CDMA primarily because of the de-multiplexing and multiplexing required for the second technique and because of the configurable multi-rate filters required for the spectrum partitioning in the third approach. This new invention implements the second and third approaches without their disadvantages and moreover provides the added performance improvements that will be realized with the use of the complex Walsh and hybrid complex Walsh codes in place of the real Walsh codes. These new 4-phase complex Walsh orthogonal CDMA codes replacing the current 2-phase real Walsh codes will provide improvements that include an increase in the carrier-to-noise ratio (CNR) for data symbol recovery in the receiver, lower correlation side-lobes under timing offsets both with and without PN spreading, lower levels of harmonic interference caused by non-linear amplification of multi-carrier CDMA signals, and reduced phase tracking jitter for code tracking to support both acquisition and synchronization. These potential performance improvements simply reflect the widely known principle that complex CDMA is better than real CDMA. The hybrid complex Walsh offers these same improvements together with the flexibility of more choices in the code lengths at the expense of increasing the number of code phases on the unit circle thereby introducing multiplications into the encoding and decoding implementations.

The new complex Walsh and hybrid complex Walsh CDMA orthogonal codes disclosed in [6] have been invented to be the

natural development for the Walsh codes and therefore are the correct complex Walsh codes to within arbitrary factors that include scale and rotation, which are not relevant to performance. This natural development of the complex Walsh codes in the N-dimensional complex code space  $C^N$  extended the correspondences between the real Walsh codes and the Fourier codes in the N-dimensional real code space  $R^N$ , to correspondences between the complex Walsh codes and the discrete Fourier transform (DFT) codes in  $C^N$ .

The new hybrid complex Walsh orthogonal CDMA codes increase the choices for the code length by allowing the combined use of complex Walsh and discrete Fourier transform complex orthogonal codes using a Kronecker construction, direct sum construction, as well as the possibility for more general functional combining.

**Transmitter equations (3)** describe a representative complex Walsh CDMA encoding for multiple data rate users for the transmitter in FIG. 1 using the definition of the complex Walsh CDMA codes in the invention application [6]. Lowest data rate users are assumed to communicate at the lowest symbol rate equal to the code repetition rate of the N chip complex Walsh code, which means they are assigned N chip code vectors from the  $N \times N$  complex Walsh code matrix  $\tilde{W}_N$  in 36 for their channelization codes. Higher data rate users will use shorter complex Walsh codes. Reference complex Walsh code matrix  $\tilde{W}_N$  has N Walsh row code vectors  $\tilde{W}_N(c)$  each of length N chips and indexed by  $c=0,1,\dots,N-1$ , with  $\tilde{W}_N(c)=[\tilde{W}_N(c,0),\dots,\tilde{W}_N(c,N-1)]$  wherein  $\tilde{W}_N(c,n)$  is chip n of code c with the possible values  $\tilde{W}_N(c,n)=+/-1 +/-j$ . Complex Walsh code vectors in the N dimensional complex code space  $C^N$  are defined using the real Walsh code vectors from the N dimensional real code space  $R^N$  for the real and complex code vectors using the equation  $\tilde{W}_N(c)=W(cr)+jW(ci)$  where the mapping of the complex Walsh code index c into the real Walsh code

indices  $cr$  and  $ci$  is defined by the mapping of  $c$  into  $cr(c)$  and  
605  $ci(c)$  in 36.

The multiple data rate menu in 37 lists the possible user  
data symbol rates  $R_s$  and the number of symbols transmitted over  
each  $N$  chip reference code length. User symbol rate  $R_s = 1/N(m)T$   
for the users in group  $m$  is equal to the number of user data  
610 symbols  $N/N(m)$  over the  $N$  chip code block multiplied by the  
symbol rate rate  $1/NT$  of the  $N$  chip code. User data rate  $R_b$  in  
bits/second is equal to  $R_b = R_s b_s$  where  $b_s$  is the number of data  
bits encoded in each data symbol. Assuming a constant  $b_s$  for all  
of the multiple data rate users, the user data rate becomes  
615 directly proportional to the user symbol rate  $R_b \sim R_s$  which means  
the user symbol rate menu in 37 is equivalent to the user data  
rate menu.

Data symbol vector 38 stores the  $N$  data symbols  
 $\{Z(u_{m,k_m})\}$  for the  $N$  chip code block in an  $1 \times N$  dimensional data  
620 symbol vector indexed by  $d = d_0 + d_1 2 + d_2 4 + \dots + d_{M-2} N/4 + d_{M-1} N/2 = 0, 1, 2, \dots, N-1$ ,  
where the binary word representation is  $d = d_0 \dots d_{M-1}$  and the  $\{d_m\}$   
are the binary coefficients. With the availability of this  $1 \times N$   
dimensional data symbol vector, it is observed that the real  
Walsh implementation for the multiple data rate users in 3 in  
625 equations (3) must assign the 2 chip data symbols  $Z(u_0, z_k)$  to the  
 $d_{M-1}$  field, the 4 chip data symbols  $Z(u_1, z_k)$  to the  $d_{M-1} d_{M-2}$  field,  
 $\dots$ , and the  $N$  chip data symbols  $Z(u_{M-1}, z_k)$  to the  $d_0 \dots d_{M-1}$  field  
in order to provide orthogonality between the code vectors in the  
different groups. For the complex Walsh the same data assignment  
630 is used with the modification that the  $N/N(m)$  data symbols for  
the  $N(m)$  chip code vectors of group  $m$  assigned to data field  
 $d_{M-m} d_{M-m+1} \dots d_{M-1}$  of  $d$  using the real Walsh, are now mapped  
into  $N/N(m)$   $N$ -chip code vectors over the same group  $m$  data field  
 $d_{M-m} d_{M-m+1} \dots d_{M-1}$  of  $d$ . This allows a fast algorithm to be used  
635 and uses the  $N$  chip codes over the  $d_{M-m} d_{M-m+1} \dots d_{M-1}$  field of  $d$   
which field occupies the same sequency band as the frequency band

for FDM. This removes the disadvantages of using technique "B" and "C" for W-CDMA, and helps to make the complex Walsh the preferred choice compared to technique "A" which is the current art preferred choice with real Walsh.

The new invention has found a means to use the same data fields of the current W-CDMA for real Walsh, for application to the complex Walsh with the added advantages of a fast transform, simultaneous transmission of the user data symbols, and the assignment of these user data symbols to a contiguous sequency band specified by the data field of d for additional isolation between users. For a fully loaded CDMA communications frequency band the N data symbols for the multiple rate users occupy the N available data symbol locations in the data symbol vector  $d = d_0 \dots d_{M-1}$ . The construction of the data symbol vector is part of this invention disclosure and provides a means for the implementation of a fast complex Walsh encoding and decoding of the multiple data rate complex Walsh CDMA. Examples 1 and 2 in 39 and 40 illustrate representative user assignments to the data fields of the data symbol vector. This mapping of the user data symbols into the data symbol vector is equivalent to setting  $c=d$  which makes it possible to develop the fast encoding algorithm 41.

**New multiple data rate complex Walsh encoding (3)  
for transmitter**

36 N chip complex Walsh code block

$\tilde{W}_N$  = complex Walsh NxN orthogonal code matrix  
consisting of N rows of N chip code vectors  
= [  $\tilde{W}_N(c)$  ] matrix of row vectors  $\tilde{W}_N(c)$

= [  $\tilde{W}_N(c,n)$  ] matrix of elements  $\hat{W}_N(c,n)$

$\tilde{W}_N(c)$  = complex Walsh code vector c  
=  $W_N(cr) + jW_N(ci)$  for  $c=0,1,\dots,N-1$

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 $W_N(cr), W_N(ci) = \text{Real Walsh } 1 \times N \text{ code vectors } cr, ci$ 
 $c = 0, 1, 2, \dots, N-1$ 
 $= \text{Real Walsh code index for } N \text{ chip block}$ 
 $= (cr, ci) \text{ Pair of real Walsh code vectors}$ 
 $cr=cr(c) \text{ and } ci=ci(c) \text{ which are assigned to}$ 

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 $\text{the real and to the imaginary axes}$ 
 $n = 0, 1, 2, \dots, N-1$ 
 $= \text{Chip index for } N \text{ chip block}$ 

Mapping of real Walsh to complex Walsh

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Complex Walsh code	Real Axis real Walsh codes	Complex Axis real Walsh codes
c	cr(c)	ci(c)
0	0	0
1, 2, ..., N/2-1	2c	2c-1
N/2	N-1	N-1
N/2+1, ..., N-1	2N-2c-1	2N-2c

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 $\tilde{W}(c, n) = \text{complex Walsh code u chip } n$ 
 $= +/-1 +/-j \text{ possible values}$ 

$$= (-1)^{[cr_{M-1}n_0 + \sum_{i=1}^{M-1} (cr_{M-1-i} + cr_{M-i})n_i]} + j(-1)^{[ci_{M-1}n_0 + \sum_{i=1}^{M-1} (ci_{M-1-i} + ci_{M-i})n_i]}$$

$$cr = \sum_{i=0}^{M-1} cr_i 2^i \text{ binary representation of } cr$$

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$$ci = \sum_{i=0}^{M-1} ci_i 2^i \text{ binary representation of } ci$$

$$n = \sum_{i=0}^{M-1} n_i 2^i \text{ binary representation of } n$$

$$= W_N(cr) + jW_N(ci) \quad \text{for } c_m=0, 1, \dots, 2^m-1$$

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### 37 Multiple data rate menu

Symbol rate menu for multiple data rates

Symbol rate	Symbol rate, Symbols/second	Symbols per N chips
$R_s =$	$1/2T$	$N/2$
$=$	$1/4T$	$N/4$
$=$	$1/8T$	$N/8$
	$\vdots$	$\vdots$
$=$	$1/2NT$	$2$
$=$	$1/NT$	$1$

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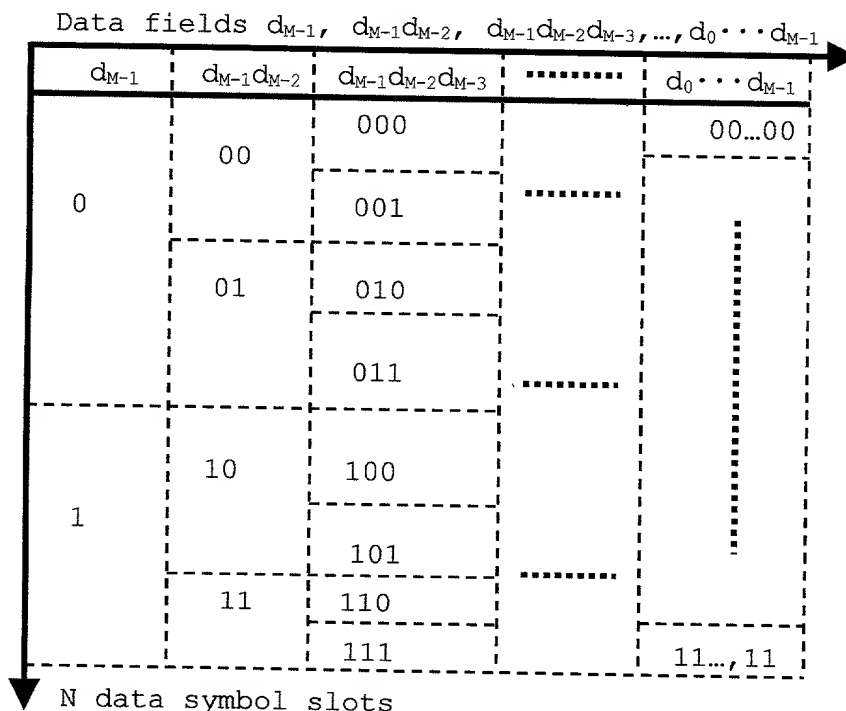
- 38 Data symbol vector field indexed by  $d=d_0+d_12+d_24+\dots+d_{M-2}N/4+d_{M-1}N/2$  is partitioned into M data fields with each assigned to one group of multiple data rate users. Writing d as a binary word  $d=d_0d_1\dots d_{M-1}$  enables the data fields to be identified as  $d_{M-1}$ ,  $d_{M-1}d_{M-2}$ ,  $d_{M-1}d_{M-2}d_{M-3}$ , ...,  $d_0\dots d_{M-1}$  which respectively are assigned to the user groups  $u_0, u_1, \dots, u_{M-1}$ .

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Menu of user assignments to the data vector fields		
User group	Available channalization codes	Field assignment in data vector
$u_0$	$u_0=0$ $u_0=1$	$d_{M-1} = 0$ $= 1$
$u_1$	$u_1=0$ $u_1=1$ $u_1=2$ $u_1=3$	$d_{M-1}d_{M-2} = 00$ $= 01$ $= 10$ $= 11$
$\vdots$	$\vdots$	$\vdots$
$u_{M-1}$	$u_{M-1}=0$ $\vdots$ $u_{M-1}=N-1$	$d_0 \cdots d_{M-1} = 00 \cdots 00$ $\vdots$ $= 11 \cdots 11$

### 39 Example 1 of multiple data rate menu:

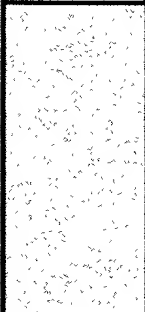
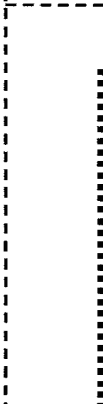



There is 1 user for each group  $u_0, u_1, \dots, u_{M-2}$  and 2 users for  $u_{M-1}$  with each user selecting the lowest sequency channel corresponding to the lowest index of channels available to the group.

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Example 1 of multiple data rate menu

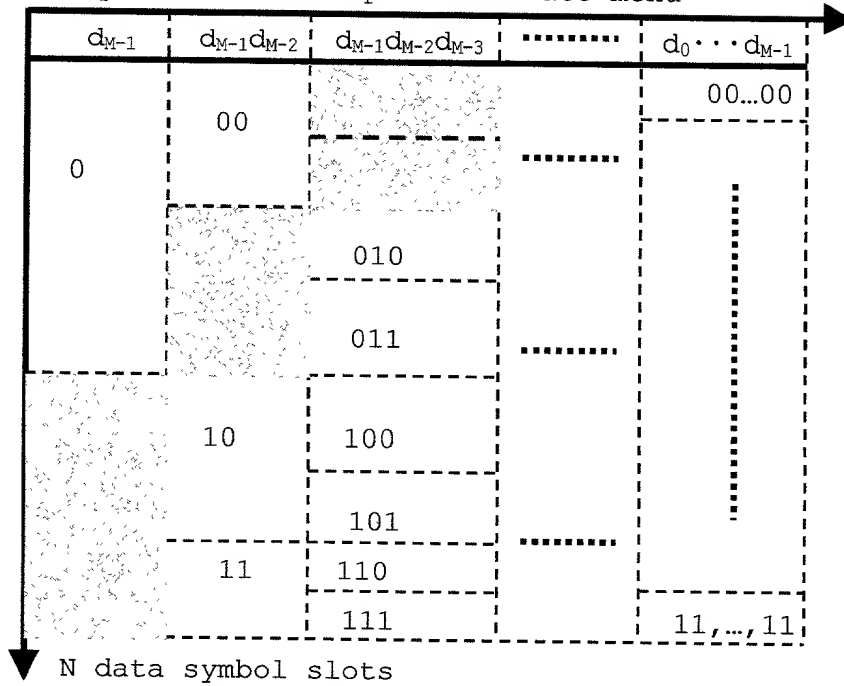
$d_{M-1}$	$d_{M-1}d_{M-2}$	$d_{M-1}d_{M-2}d_{M-3}$	.....	$d_0 \cdots d_{M-1}$
	00	000		00...00
	01	001	.....	
		010		
		011	.....	
1		100		
		101		
	11		.....	
		111		
N data symbol slots				

40 Example 2 of multiple data rate menu:

There is 1 user in each group  $u_0$  and  $u_1$  and 2 users in  $u_2$  with each user selecting the highest sequency channel corresponding to the highest index of channels available to the group.



Example 2 of multiple data rate menu



- 41 Complex Walsh encoding and channel combining uses a computationally efficient fast encoding algorithm. This algorithm implements the encoding with an  $M$  pass computation. Passes  $1, 2, 3, \dots, M$  respectively perform the  $2, 4, 8, \dots, N$  chip complex Walsh encoding of the data symbol vector successively starting with the 2 chip encoding in pass 1, the 4 chip encoding in passes  $1, 2$ , the 8 chip encoding in passes  $1, 2, 3$ , and the  $N$  chip encoding in passes  $1, 2, 3, \dots, M$  where  $N=2^M$ . Using the binary word representations for both  $d$  and  $n$ , this  $M$  pass algorithm is:

$$\begin{aligned} \text{Pass 1: } & Z^{(1)}(n_{M-1}d_1 \dots d_{M-1}) \\ &= \sum_{d_0=dr_0=di_0=0,1} Z(d_0 \dots d_{M-1}) [(-1)^{dr_0 n_{M-1}} + j(-1)^{di_0 n_{M-1}}] \end{aligned}$$

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Pass m for m=2,...,M-1

$$Z^{(m)}(n_{M-1} \cdots n_{M-m} d_m \cdots d_{M-1})$$

$$= \sum Z^{(m-1)}(n_{M-1} \cdots n_{M-m+1} d_{m-1} \cdots d_{M-1}) \cdot$$

$$\uparrow [(1)^{dr_{m-1}(n_{M-m}+n_{M-m+1})+j} (-1)^{di_{m-1}(n_{M-m}+n_{M-m+1})}]$$

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$$d_{m-1}=dr_{m-1}=di_{m-1}=0,1$$

Pass M:  $Z^{(M)}(n_{M-1}n_{M-2} \cdots n_1n_0)$

$$= \sum Z^{(M-1)}(n_{M-1}n_{M-2} \cdots n_1 d_{M-1}) \cdot$$

$$\uparrow [(1)^{dr_{M-1}(n_0+n_1)+j} (-1)^{di_{M-1}(n_0+n_1)}]$$

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$$d_{M-1}=dr_{M-1}=di_{M-1}=0,1$$

$$= \tilde{Z}(n_{M-1}n_{M-2} \cdots n_1n_0)$$

860

An additional re-ordering pass is added to change the encoded N chip block  $\tilde{Z}(n_{M-1}n_{M-2} \cdots n_1n_0)$  in bit reversed ordering to the normal readout ordering

$$\hat{Z}(n_0n_1 \cdots n_{M-2}n_{M-1}) = \tilde{Z}(n)$$

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## 42 PN scrambling

$P_R(n), P_I(n)$  = PN code chip n for real and  
Imaginary axes

$Z(n)$  = PN scrambled complex Walsh encoded data  
chips after summing over the users

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$$= \sum_u \tilde{Z}(n) [P_R(n) + j P_I(n)]$$

$$= \sum_u \tilde{Z}(n) [\text{sgn}\{P_R(n)\} + j \text{sgn}\{P_I(n)\}]$$



sign of the PN chip. Purpose of the PN encoding for complex data symbols is to provide scrambling of the summed complex Walsh encoded data symbols as well as isolation between groups of users. Output of this complex Walsh CDMA encoding followed by the complex PN scrambling are the CDMA encoded chips over the N chip block  $\{Z(n)\}$ .

**Transmitter equations (6) for hybrid complex Walsh orthogonal encoding** of multiple data rate users are derived by starting with the hybrid complex Walsh orthogonal codes disclosed in the invention application [6]. The discrete Fourier transform (DFT) CDMA codes used in the example generation of hybrid complex Walsh orthogonal CDMA codes in [6] are given in equations (4) along with a fast encoding algorithm.

**N-chip DFT complex orthogonal CDMA codes (4)**

**43 DFT code vectors**

$E_N$  = DFT  $N \times N$  orthogonal code matrix consisting of  
N rows of N chip code vectors

=  $[E_N(c)]$  matrix of row vectors  $E(c)$

=  $[E_N(c, n)]$  matrix of elements  $E(c, n)$

$E_N(c)$  = DFT code vector  $c$

=  $[E_N(c, 0), E_N(c, 1), \dots, E_N(c, N-1)]$

=  $1 \times N$  row vector of chips  $E_N(c, 0), \dots, E_N(c, N-1)$

$E_N(c, n)$  = DFT code  $c$  chip  $n$

=  $e^{j2\pi cn/N}$

=  $\cos(2\pi cn/N) + j\sin(2\pi cn/N)$

= N possible values on the unit circle

**44 Fast encoding algorithm for N chip block of data**  
in the data vector  $d = d_0 d_1 \dots d_{M-2} d_{M-1}$

Pass 1:  $Z^{(1)}(d_0 d_1 \dots d_{M-2} n_0)$

=  $\sum Z(d_0 d_1 d_2 \dots d_{M-2} d_{M-1}) (e^{j2\pi})^{d_{M-1} n_0 / 2}$

↑

$d_{M-1} = 0, 1$

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Pass m for m=2,...,M-1

$$\begin{aligned}
 & Z^{(m)}(d_0 \cdots d_{M-m-1} n_{m-1} \cdots n_0) \\
 &= \sum_{d_{M-m}} Z^{(m-1)}(d_0 \cdots d_{M-m} n_{m-2} \cdots n_0) \bullet \\
 & \quad e^{j2\pi [d_{M-m}(n_0 + n_1 2 + \cdots + n_{m-1} 2^{(m-1)}) / 2^m]}
 \end{aligned}$$

950

Pass M for m=2,...,M-1

$$\begin{aligned}
 & Z^{(M)}(n_{M-1} \cdots n_0) \\
 &= \sum_{d_0} Z^{(m-1)}(d_0 n_{M-2} \cdots n_0) \bullet \\
 & \quad e^{j2\pi [d_0(n_0 + n_1 2 + \cdots + n_{M-1} 2^{(M-1)}) / 2^M]} \\
 &= \tilde{Z}(n_{M-1} n_{M-2} \cdots n_1 n_0)
 \end{aligned}$$

960

An additional re-ordering pass is added to change the encoded N chip block  $\tilde{Z}(n_{M-1} n_{M-2} \cdots n_1 n_0)$  in bit reversed ordering to the normal readout ordering

$$\tilde{Z}(n_0 n_1 \cdots n_{M-2} n_{M-1}) = \tilde{Z}(n)$$

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An additional re-ordering pass is added to change the encoded N chip block  $\tilde{Z}(n_{M-1} n_{M-2} \cdots n_1 n_0)$  in bit reversed ordering to the normal readout ordering

$$\tilde{Z}(n_0 n_1 \cdots n_{M-2} n_{M-1}) = \tilde{Z}(n)$$

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DFT code matrix and the row code vectors are defined in **43** for an N chip block. A fast algorithm for the encoding of the N chip data vector  $Z(d_0 d_1 d_2 \cdots d_{M-2} d_{M-1})$  is defined in **44** in a format similar to the fast algorithm for the complex Walsh encoding in equations (3). It is well known that the computational complexity of the fast DFT encoding algorithm is  $R_A \approx 2M$  real

additions per data symbol plus  $R_M \approx 3M$  real multiplications per  
975 data symbol. The relatively high complexity implementation cost  
of multiplies makes it desirable to limit the use of DFT codes to  
applications such as the hybrid complex Walsh wherein the number  
of real multiplies per data symbol can be kept more reasonable.

Hybrid complex Walsh orthogonal CDMA codes increase  
980 the flexibility in choosing the code lengths for multiple data  
rate users at the implementation cost of introducing multiply  
operations into the CDMA encoding and decoding. Two of several  
means for construction given in the patent application [6] are  
985 the Kronecker product and the direct sum. The direct sum will not  
be considered since the addition of the zero matrix in the  
construction is generally not desirable for CDMA communications  
although the direct sum construction provides greater flexibility  
in the choice of  $N$  without necessarily introducing a multiply  
penalty. Using the Kronecker product construction in reference  
990 [6] the hybrid complex Walsh orthogonal CDMA codes can be  
constructed as demonstrated in equations (4).

Equations (5) list construction and examples of the hybrid  
complex Walsh orthogonal CDMA codes using the Kronecker approach  
and DFT matrices  $E_N$  to expand the complex Walsh to a hybrid  
995 complex Walsh. Low order CDMA code examples 45 illustrate  
fundamental relationships between the DFT, complex Walsh, and the  
real Walsh or equivalently Hadamard. Kronecker construction is  
defined in 46. CDMA current and developing standards use the  
prime 2 which generates a code length  $N=2^M$  where  $M$ =integer. For  
1000 applications requiring greater flexibility in code length  $N$ ,  
additional primes can be used using the Kronecker construction. We  
illustrate this in the examples 47 with the addition of  
prime=3. The use of prime=3 in addition to the prime=2 in the  
range of  $N=8$  to 64 is observed to increase the number of  $N$   
1005 choices from 4 to 9 at a modest cost penalty of using multiples  
of the angle increment 30 degrees for prime=3 in addition to the  
angle increment 90 degrees for prime=2. As noted in 46 there

are several choices in the ordering of the Kronecker product construction and 2 of these choices are used in the construction.

1010 **Examples of hybrid complex Walsh orthogonal codes (5)**

45 Examples of low-order codes

$$\begin{aligned} 2 \times 2 \quad E_2 &= \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \\ &= (e^{-j\pi/4} / \sqrt{2}) * \tilde{W}_2 \\ &= H_2 \quad 2 \times 2 \text{ Hadamard} \end{aligned}$$

$$3 \times 3 \quad E_3 = \begin{bmatrix} 1 & 1 & 1 \\ 1 & e^{j2\pi/3} & e^{j2\pi/3} \\ 1 & e^{j2\pi/3} & e^{j2\pi/3} \end{bmatrix}$$

$$4 \times 4 \quad \tilde{W}_4 = \begin{bmatrix} 1+j & 1+j & 1+j & 1+j \\ 1+j & -1+j & -1-j & 1-j \\ 1+j & -1-j & 1+j & -1-j \\ 1+j & 1-j & -1-j & -1+j \end{bmatrix}$$

$$E_4 = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & j & -1 & -j \\ 1 & -1 & 1 & -1 \\ 1 & -j & -1 & j \end{bmatrix}$$

$$= (e^{-j\pi/4} / \sqrt{2}) \tilde{W}_4$$

46 Kronecker product construction for  $N = \prod_k N_k$

Code matrix  $C_N = N \times N$  hybrid orthogonal CDMA code matrix  
Kronecker product construction of  $C_N$

$$C_N = C_0 \prod_{k>0} \otimes C_{N_k}$$

Kronecker product definition

A =  $N_a \times N_a$  orthogonal code matrix

B =  $N_b \times N_b$  orthogonal code matrix

$A \otimes B$  = Kronecker product of matrix A and matrix B

1045 =  $N_a N_b \times N_a N_b$  orthogonal code matrix consisting  
of the elements  $[a_{ik}]$  of matrix A multiplied  
by the matrix B

=  $[a_{ik} B]$

47 Kronecker product construction examples for primes

$p=2,3$  and the range of sizes  $8 \leq N \leq 64$

8x8  $C_8 = \tilde{W}_8$

12x12  $C_{12} = \tilde{W}_4 \otimes E_3$

$C_{12} = E_3 \otimes \tilde{W}_4$

16x16  $C_{16} = \tilde{W}_{16}$

18x18  $C_{18} = \tilde{W}_2 \otimes E_3 \otimes E_3$

$C_{18} = E_3 \otimes E_3 \otimes \tilde{W}_2$

24x24  $C_{24} = \tilde{W}_8 \otimes E_3$

$C_{24} = E_3 \otimes \tilde{W}_8$

32x32  $C_{32} = \tilde{W}_{32}$

36x36  $C_{36} = \tilde{W}_4 \otimes \tilde{W}_3 \otimes \tilde{W}_3$

$C_{36} = \tilde{W}_3 \otimes \tilde{W}_3 \otimes \tilde{W}_4$

48x48  $C_{48} = \tilde{W}_{16} \otimes \tilde{W}_3$

$C_{48} = \tilde{W}_3 \otimes \tilde{W}_{16}$

64x64  $C_{64} = \tilde{W}_{64}$

1065 A fast algorithm for the encoding of the hybrid complex  
Walsh CDMA orthogonal codes is described in equations (6) for the  
representative example 48 which constructs the  $N \times N$  hybrid  
complex Walsh orthogonal CDMA code matrix  $C_N = \tilde{W}_{N_0} \otimes E_{N_1}$  as the



Kronecker product of the  $N_0 \times N_0$  complex Walsh  $\tilde{W}_{N_0}$  and the  $N_1 \times N_1$   
 complex DFT, where  $N = N_0 N_1$ . Each chip element of  $C_N$  is the product  
 of the chip elements of the complex Walsh and complex DFT  
 code matrices. The complex Walsh and DFT codes are phase codes  
 which means the phase of each  $C_N$  chip element is the sum of the  
 phases of the chip elements for the complex Walsh and complex  
 DFT. Chip element equations are  $C_N(c, n) = \tilde{W}_{N_0}(c\tilde{w}, n\tilde{w}) E_{N_1}(ce, ne)$  with  
 $c = ce + c\tilde{w} N_1$  and  $n = ne + n\tilde{w} N_1$ . For multiple data rate  
 data symbol assignments and for the construction of the fast  
 encoding algorithm, it is convenient to use a binary word  
 representation of the chip element indices  $c, n$ . Binary word  
 representation is  $c = ce_0 ce_1 \dots ce_{M_1-1} c\tilde{w}_{M_1} c\tilde{w}_{M_1+1} \dots c\tilde{w}_{M-1} = c_0 c_1 \dots c_{M-2} c_{M-1}$   
 where the first binary word is a function of the binary words  
 for the complex Walsh and complex DFT code indices, and the  
 second binary word is a direct representation of the  $C_N$  indices  
 which will be used for the data vector construction. The same  
 binary word representations apply for the chip index  $n$  upon  
 substituting the  $n$  for  $c$ . User data in equations (3) for  
 the  $N$  chip code block is mapped into the  $N$  data symbol vector  $d =$   
 $d_0 \dots d_{M-1}$  which is obtained from the binary word for  $c$  by  
 substituting the index  $d$  for the index  $c$  in the binary word  
 representation.

The multiple data rate data symbol mapping in equations  
 (6) for the hybrid complex Walsh codes remains the same as used  
 in 38, 39, 40 in equations (3) for the complex Walsh codes. The  
 data symbol mapping assigns the  $N/2$  data symbols of the 2 chip  
 data symbol transmission rate users to the  $d_{M-1}$  field, the  $N/4$   
 data symbols of the 4 chip data symbol transmission rate users  
 are assigned to the  $d_{M-1} d_{M-2}$  field, ..., and the single data  
 symbols of the  $N$  chip data symbol transmission rate users are  
 assigned to the  $d_0 \dots d_{M-1}$  field, where the data vector index " $d$ "  
 is represented as the binary number  $d = d_0 \dots d_{M-1}$  and the  $\{d_m\}$  are  
 the binary coefficients. For a fully loaded CDMA communications

frequency band the N data symbols occupy the N available data symbol locations in the data symbol vector  $d = d_0 \cdots d_{M-1}$ . The menu of available user assignments to the data vector fields is given in 38 in equations (3). Examples 1 and 2 in 39 and 40 in equations (3) illustrate representative user assignments to the data fields of the data symbol vector. This mapping of the user data symbols into the data symbol vector is equivalent to setting  $c=d$  which makes it possible to develop the fast encoding algorithm 51.

**Fast multiple data rate hybrid complex Walsh encoding (6) for transmitter**

48 The fast algorithm will be described for the example  $N \times N$  complex orthogonal CDMA code matrix  $C_N$  which is generated by the Kronecker product of the  $N_0 \times N_0$  complex Walsh matrix  $\tilde{W}_{N_0}$  and the complex  $N_1 \times N_1$  DFT matrix  $E_{N_1}$

$$C_N = \text{Kronecker product of } \tilde{W}_{N_0} \text{ and } E_{N_1}$$

$$= \tilde{W}_{N_0} \otimes E_{N_1}$$

$$\text{where } N = N_0 N_1$$

$$= 2^M$$

$$M = M_0 + M_1$$

$$N_0 = 2^{M_0}$$

$$N_1 = 2^{M_1}$$

49 N chip hybrid complex Walsh code block  $C_N$

$C_N$  = hybrid complex Walsh  $N \times N$  orthogonal code matrix consisting of N rows of N chip code vectors

$$= [C_N(c)] \text{ matrix of row vectors } C_N(c)$$

$$= [C_N(c,n)] \text{ matrix of elements } C_N(c,n)$$

$$C_N(c,n) = \text{hybrid complex Walsh code } c \text{ chip } n$$

$$= \tilde{W}_{N_0}(c\tilde{w}, n\tilde{w}) E_{N_1}(ce, ne)$$

$$= [+/-1 +/-j] E_{N_1}(ce, ne) \text{ values}$$

$$\text{where } c = ce + c\tilde{w} N_1$$

$$n = ne + n\tilde{w} N_1$$

50 Binary indexing of codes in the matrix  $C_N$

$$\begin{aligned}
 c &= ce_0 + ce_1 2 + \dots + ce_{M_1-1} 2^{(M_1-1)} \\
 &\quad + c\tilde{w}_{M_1} 2^{M_1} + c\tilde{w}_{M_1+1} 2^{(M_1+1)} + \dots + c\tilde{w}_{M-1} 2^{M-1} \\
 &= ce_0 ce_1 \dots ce_{M_1-1} c\tilde{w}_{M_1} c\tilde{w}_{M_1+1} \dots c\tilde{w}_{M-1} \quad \text{Binary word} \\
 n &= ne_0 + ne_1 2 + \dots + ne_{M_1-1} 2^{(M_1-1)} \\
 &\quad + n\tilde{w}_{M_1} 2^{M_1} + n\tilde{w}_{M_1+1} 2^{(M_1+1)} + \dots + n\tilde{w}_{M-1} 2^{M-1} \\
 &= ne_0 ne_1 \dots ne_{M_1-1} n\tilde{w}_{M_1} n\tilde{w}_{M_1+1} \dots n\tilde{w}_{M-1} \quad \text{Binary word}
 \end{aligned}$$

51 The fast encoding algorithm starts with the data symbol vector  $d$  and mapping of the user groups  $u_0, u_1, \dots, u_{M-1}$  into the data fields of  $d$ . This mapping is identical to the mapping defined in equations (3) for the multiple data rate complex Walsh orthogonal encoding of the CDMA over an  $N$  chip block. However, the fast algorithm for the hybrid complex Walsh encoding is modified to accommodate the Kronecker construction as illustrated by the following fast algorithm for the hybrid complex Walsh example in 48. Using the binary representations of  $d, n$

$$\begin{aligned}
 d &= d_0 d_1 \dots d_{M_1-1} d_{M_1} \dots d_{M-1} \\
 &= de_0 de_1 \dots de_{M_1-1} d\tilde{w}_{M_1} \dots d\tilde{w}_{M-1} \\
 n &= n_0 n_1 \dots n_{M_1-1} n_{M_1} \dots n_{M-1} \\
 &= ne_0 ne_1 \dots ne_{M_1-1} n\tilde{w}_{M_1} \dots n\tilde{w}_{M-1}
 \end{aligned}$$

and the same approach used to derive the fast algorithms 41 in equations (3) and 44 in equations (4), enables the  $M$  pass fast algorithm to be defined

Pass 1 for complex Walsh codes

$$\begin{aligned}
 &Z^{(1)}(d_0 \dots d_{M_1-1} n_{M_1} d_{M_1+1} \dots d_{M-1}) \\
 &= \sum Z(d_0 \dots d_{M-1}) \cdot
 \end{aligned}$$



$$[e^{j2\pi} d_0 (n_{M_0} + n_{M_0+1}^2 + \dots + n_{M-1}^{2^{M_0-1}}) / 2^{M_0}]$$

$$d_0 = 0, 1$$

$$= \tilde{Z}(n_{M-1}n_{M-2} \dots n_1n_0)$$

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An additional re-ordering pass is added to change the encoded N chip block  $\tilde{Z}(n_{M-1}n_{M-2} \dots n_1n_0)$  in bit reversed ordering to the normal readout ordering

$$\hat{Z}(n_0n_1 \dots n_{M-2}n_{M-1}) = \tilde{Z}(n)$$

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The fast algorithm in 51 is a computationally efficient means to implement the hybrid complex Walsh encoding of each N chip code block for multiple data rate users whose lowest data rate corresponds to the data symbol rate of an N chip encoded user. The computational complexity of this fast encoding algorithm can be estimated using the computational complexities of the complex Walsh and the DFT fast encoding algorithms, which gives the estimate:  $R_A \approx 2M + M_1 + 2$  real additions per data symbol, and  $R_M \approx 2M_1$  real multiplies per data symbol.

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The fast algorithm in 51 consists of M signal processing passes on the stored data symbols, followed by a re-ordering pass for readout of the N chip block of encoded data symbols. Advantage is taken of the equality  $c=d$  which allows the d to be used in the code indices for the complex Walsh:  $d_m = c_m$ ,  $d_r = c_r$ ,  $d_i = c_i$ . Pass 1 implements 2 chip encoding, passes  $m=2, \dots, M_0$  implement  $2^m$  chip encoding with the complex Walsh codes, passes  $M_0+1, M_0+2, \dots, M_0+M_1-1 = M-1$  implement  $2^{M_0+m}$  chip encoding with the complex DFT codes, and the last pass M encodes the  $N=2^M$  chip data symbols with the DFT codes. This fast algorithm only differs from the fast algorithm in 46 in equations (4) in the use of both the complex Walsh codes and the complex DFT codes with their Kronecker indexing. Unlike the fast algorithm for the real Walsh encoding as well as the algorithm for the complex DFT encoding,

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the complex Walsh portion of the fast algorithm 51 uses both the sign of the complex Walsh code from the current pass and from the previous pass starting with pass 2.

The generalization of the fast algorithm in 51 in equations (6) to other Kronecker product constructions for  $C_N$  and to the more general constructions for  $C_N$  discussed in reference [6] should be apparent to anyone skilled in the CDMA communications art.

**Receiver equations (7)** describe a representative multiple data rate complex Walsh CDMA decoding for multiple data users for the receiver in FIG. 3 using the definition of the complex Walsh CDMA codes in the invention application [6]. The receiver front end 52 provides estimates  $\{\hat{Z}(n)\}$  of the transmitted multiple data rate complex Walsh CDMA encoded chips  $\{Z(n)\}$ . Orthogonality property 53 is expressed as a matrix product of the complex Walsh code chips or equivalently as a matrix product of the complex Walsh code chip numerical signs of the real and imaginary components, for any of the 2, 4, 8, ..., N/2, N chip complex Walsh channelization codes and their repetitions over the N chip code block. The 2-phase PN codes 54 have the useful decoding property that the square of each code chip is unity which is equivalent to observing that the square of each code chip numerical sign is unity.

**Receiver decoding of complex Walsh and hybrid complex Walsh CDMA encoded chips (7)**

52 Receiver front end in FIG. 3 provides estimates

$\{\hat{Z}(n)\}$  28 of the encoded transmitter chip symbols  $\{Z(n)\}$  41 in equations (3)

53 Orthogonality properties of the complex Walsh  $N \times N$  matrix

$$\sum_n \tilde{W}_N(\hat{c}, n) \tilde{W}_N^*(n, c) = \sum_n [\text{sgn}\{W_N(\hat{c}r, n) + j \text{sgn}\{W_N(\hat{c}i, n)\}\}][\text{sgn}\{W_N(n, cr)\} - j \text{sgn}\{W_N(n, ci)\}]$$

$$1250 \quad \quad \quad = 2N \delta(\hat{c}, c)$$

where  $\hat{c}, c, n = 0, 1, \dots, N$

$$\begin{aligned} \delta(\hat{c}, c) &= \text{Delta function of } \hat{c} \text{ and } c \\ &= 1 \quad \text{for } \hat{c} = c \\ &= 0 \quad \text{otherwise} \end{aligned}$$

1255  $cr=cr(c), ci=ci(c)$  are defined  
in equations (3)

54 PN de-scrambling of the receiver estimates of the complex and hybrid complex Walsh encoded data chips

1260  $P_R(n), P_I(n) =$  PN code chip  $n$  for real and imaginary axes

$$\begin{aligned} \tilde{Z}(n) &= \text{PN de-scrambled receiver estimates of the} \\ &\quad \text{transmitted CDMA encoded chips } \hat{Z}(n) \\ &= \hat{Z}(n) [P_R(n) - jP_I(n)] \end{aligned}$$

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55a Complex Walsh decoding and uses a computationally efficient fast encoding algorithm. This algorithm implements the decoding with an  $M$  pass computation. Passes  $1, 2, 3, \dots, M$  respectively perform the  $2, 4, 8, \dots, N$  chip complex Walsh decoding of the data symbol vector successively starting with the 2 chip decoding in pass 1, the 4 chip decoding in passes 1, 2, and the  $N$  chip decoding in passes  $1, 2, 3, \dots, M$  where  $N=2^M$ . Using the binary word representations for both  $d$  and  $n$ , this  $M$  pass algorithm is:

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Pass 1:

$$\begin{aligned} \hat{Z}^{(1)}(d_{M-1}n_1n_2 \cdots n_{M-2} n_{M-1}) \\ = \sum_{\substack{\uparrow \\ n_0=0,1}} \tilde{Z}(n_0n_1n_2 \cdots n_{M-2}n_{M-1}) \bullet \\ \quad \quad \quad [(-1)^{n_0d_{M-1}} - j(-1)^{n_0d_{M-1}}] \end{aligned}$$

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Pass m for m=2,...,M-1

$$\begin{aligned} & \hat{Z}^{(m)}(d_{M-1}d_{M-2}\cdots d_{M-m}n_m\cdots n_{M-2}n_{M-1}) \\ &= \sum \hat{Z}^{(m-1)}(d_{M-1}d_{M-2}\cdots d_{M-m+1}n_{m-1}\cdots n_{M-2}n_{M-1}) \cdot \\ & \quad \uparrow [(-1)^{n_{m-1}}(dr_{M-m}+dr_{M-m+1}) - j(-1)^{n_{m-1}}(di_{M-m}+di_{M-m+1})] \\ & n_{m-1} = 0, 1 \end{aligned}$$

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Pass M

$$\begin{aligned} & \hat{Z}^{(M)}(d_{M-1}d_{M-2}\cdots d_0) \\ &= \sum \hat{Z}^{(M-1)}(d_{M-1}d_{M-2}\cdots d_1n_0) \cdot \\ & \quad \uparrow [(-1)^{n_{M-1}}(dr_0+dr_1) - j(-1)^{n_{M-1}}(di_0+di_1)] \\ & n_{M-1} = 0, 1 \\ &= N \hat{Z}(d_{M-1}d_{M-2}\cdots d_0) \end{aligned}$$

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An additional re-ordering pass is added to change the decoded N chip block  $\hat{Z}(d_{M-1}d_{M-2}\cdots d_0)$  in bit reversed ordering to the normal readout ordering. In this representative fast implementation the scaling factor N has been removed in the re-ordering pass whereas a typical implementation will re-scale each pass. The output of this final pass is the receiver estimate of the transmitted data symbol vector

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$$\hat{Z}(d_0d_2\cdots d_{M-1}) = \hat{Z}(d)$$

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**55b** Hybrid complex Walsh decoding uses a computationally efficient fast encoding algorithm. Similar to the complex Walsh this algorithm implements the decoding with an M pass computation 1,2,3,...,M:



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Pass  $m=1, \dots, M_1$  for complex DFT codes

$$\hat{Z}^{(m)}(n_0 \dots n_{M-m-1} d_{m-1} \dots d_0)$$

$$= \sum \hat{Z}^{(m-1)}(n_0 \dots n_{M-m} d_{m-2} \dots d_0) \bullet$$

$$[e^{j2\pi} n_{M-m} (d_0 + d_1 2 + \dots + d_{m-1} 2^{m-1}) / 2^m]$$



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$$n_{M-m} = 0, 1$$

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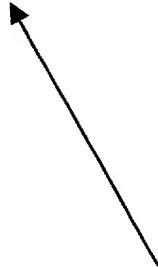
Pass  $M_1+m=M_1+1, M_1+2, \dots, M-1$  for complex Walsh codes

$$\hat{Z}^{(M_1+m)}(d_{M-1} \dots d_{M-m} n_m \dots n_{M_0-1} d_{M_1-1} \dots d_0)$$

$$= \sum \hat{Z}^{(M_1+m-1)}(d_{M-1} \dots d_{M-m-1} n_{m-1} \dots n_{M_0-1} d_{M_1-1} \dots d_0) \bullet$$

$$[(-1)^{n_{m-1}} (dr_{M_0-m} + dr_{M_0-m+1})$$

$$-j(-1)^{n_{m-1}} (di_{M_0-m} + di_{M_0-m+1})]$$



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$$n_{m-1} = 0, 1$$

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Pass M for complex Walsh codes

$$\begin{aligned}
& \hat{Z}^{(M)}(d_{M-1} \cdots d_1 d_0) \\
&= \sum_{n_{M_0-1}=0,1} \hat{Z}^{(M-1)}(d_{M-1} \cdots d_{M-m-1} n_{M_0-1} d_{M_1-1} \cdots d_0) \cdot \\
&\quad [(-1)^{n_{M_0-1}(dr_0 + dr_1)}] \\
&\quad -j(-1)^{n_{M_0-1}(di_0 + di_1)} \\
&= \hat{Z}(d_{M-1} d_{M-2} \cdots d_1 d_0)
\end{aligned}$$

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An additional re-ordering pass is added to change the decoded N chip block  $\hat{Z}(d_{M-1} d_{M-2} \cdots d_0)$  in bit reversed ordering to the normal readout ordering. In this representative fast implementation the scaling factor N has been removed in the re-ordering pass whereas a typical implementation will re-scale each pass. The output of this final pass is the receiver estimate of the transmitted data symbol vector

$$\hat{Z}(d_0 d_2 \cdots d_{M-1}) = \hat{Z}(d)$$

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The fast decoding algorithms **55a**, **55b** perform the inverse of the signal processing for the encoding **41**, **51** in equations (3), (6) of the complex, hybrid complex Walsh respectively, to recover estimates  $\{\hat{Z}(d)\}$  of the transmitter user data symbols  $\{Z(d)\}$ . These algorithms are computationally efficient means to implement the complex and hybrid complex Walsh decoding of each N chip code block for multiple data rate users whose lowest data rate corresponds to the data symbol rate of an N chip encoded

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user. For the fast complex Walsh decoding algorithm in **55a** the number of required real additions  $R_A$  per data symbol is approximately equal to  $R_A \approx 2M+2$  which is identical to the complexity metric for the fast encoding algorithm. For the fast hybrid complex Walsh decoding algorithm in **55b** the computational complexity is  $R_A \approx 2M+M_1+2$  real additions per data symbol and  $R_M \approx 2M_1$  real multiplies per data symbol which is identical to the complexity metric for the fast encoding algorithm.

For the complex Walsh decoding the fast algorithm **55a** implements  $M$  signal processing passes on the  $N$  chip block of received data chips after de-scrambling, followed by a re-ordering pass of the receiver recovered estimates of the data symbols. Passes  $m=1,2,\dots,M$  implement  $2^m$  chip decoding. For the hybrid complex Walsh the fast algorithm **55b** combines the complex Walsh algorithm with a DFT algorithm in  $M$  signal processing passes where  $M=M_0+M_1$  with  $M_0$ ,  $M_1$  respectively designating the complex Walsh, DFT decoding passes. Passes  $m=1,\dots,M_1$  implement the complex DFT decoding and the remaining passes  $M_1+1,\dots,M-1$  implement decoding with the complex Walsh codes, and the last pass  $M$  completes the complex decoding.

**FIG. 5 complex/hybrid complex Walsh CDMA encoding** is a representative implementation of the complex and hybrid complex (complex/hybrid complex) Walsh CDMA encoding which replaces the current real Walsh encoding **13** in FIG. 1, and is defined in equations (3) and (6). The input user data symbols  $\{Z(u_{m,k_m})\}$  are mapped into the data symbol vector  $Z(d)$  as described in equations (3). Data symbols  $\{Z(d)\}$  are encoded and summed over the user data symbols in **58** and **59** by the fast encoding algorithm in equations **41** in (3) for the complex Walsh and in equations **51** in (6) for the hybrid complex Walsh. For the hybrid complex Walsh, the fast complex DFT encoding **59** follows the fast complex Walsh encoding **58**. This encoding and summing over the user data symbols is followed by PN encoding with the

scrambling sequence  $[P_R(n)+jP_I(n)]$  60. Output is the stream of complex CDMA encoded chips  $\{Z(n)\}$  61.

It should be obvious to anyone skilled in the communications art that this example implementation in FIG. 5 clearly defines the fundamental CDMA signal processing relevant to this invention disclosure and it is obvious that this example is representative of the other possible signal processing approaches.

FIG. 6 complex/hybrid complex Walsh CDMA decoding is a representative implementation of complex/hybrid Walsh CDMA decoding which replaces the current real Walsh decoding 27 in FIG. 3 and is defined in equations (7). Inputs are the received estimates of the complex CDMA encoded chips  $\{\hat{Z}(n)\}$  62. The PN scrambling code is stripped off from these chips 63 by changing the sign of each chip according to the numerical sign of the real and imaginary components of the complex conjugate of the PN code as per the decoding algorithms 54 in equations (7). The complex/hybrid complex Walsh channelization coding is removed by the fast decoding algorithms in equations 55 in (7) for the complex/hybrid complex Walsh, to recover the receiver estimates  $\{\hat{Z}(d)\}$  of the transmitted data symbols  $\{Z(d)\}$ . The complex Walsh fast decoding 64 is followed by the complex DFT fast decoding 65 for the hybrid complex Walsh. Decoded outputs are the estimated data vector  $\hat{Z}(d)$  66 whose entries are read out as the set of receiver estimates  $\{\hat{Z}(u_{m,k_m})\}$  67 of the transmitted data symbols.

It should be obvious to anyone skilled in the communications art that this example implementation in FIG. 6 clearly defines the fundamental CDMA signal processing relevant to this invention disclosure and it is obvious that this example is representative of the other possible signal processing approaches.

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**Preferred embodiments** in the previous description is provided to enable any person skilled in the art to make or use the present invention. The various modifications to these embodiments will be readily apparent to those skilled in the art, and the generic principles defined herein may be applied to other embodiments without the use of the inventive faculty. Thus, the present invention is not intended to be limited to the embodiments shown herein but is not to be accorded the wider scope consistent with the principles and novel features disclosed herein.

It should be obvious to anyone skilled in the communications art that this example implementation of the complex Walsh and hybrid complex Walsh for multiple data rate users in equations (3),..., (7) clearly defines the fundamental CDMA signal processing relevant to this invention disclosure and it is obvious that this example is representative of the other possible signal processing approaches. For example, the Kronecker matrices  $E_N$  and  $H_N$  can be replaced by functionals.

For cellular applications the transmitter description which includes equations (18) describes the transmission signal processing applicable to this invention for both the hub and user terminals, and the receiver corresponding to the decoding of equations (18) describes the corresponding receiving signal processing for the hub and user terminals for applicability to this invention.

For optical communications applications the the microwave processing at the front end of both the transmitter and the receiver is replaced by the optical processing which performs the complex modulation for the optical laser transmission in the transmitter and which performs the optical laser receiving function of the microwave processing to recover the complex baseband received signal.